Chapter 8: First Order Logic

Outline

- Why FOL?
- Syntax and semantics of FOL
- Using FOL
- Wumpus world in FOL
- Knowledge engineering in FOL

Pros and cons of Propositional Logic

© Propositional logic is declarative

© Propositional logic allows partial/disjunctive/negated information

– (unlike most data structures and databases)

OPPOPOSITIONAL LOGIC IS COMPOSITIONAL:

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— meaning of $B_{1,1} \wedge P_{1,2}$ is derived from meaning of $B_{1,1}$ and of $P_{1,2}$

② Meaning in propositional logic is context-independent

- (unlike natural language, where meaning depends on context)

First-order logic

Whereas propositional logic assumes the world contains facts,

 First-order logic (like natural language) assumes the world contains

Objects: people, houses, numbers, colors, baseball games, wars, …

 Relations: red, round, prime, brother of, bigger than, part of, comes between, ...

Syntax of FOL: Basic elements

- Constants KingJohn, 2, NUS,...
- Predicates Brother, >,...
- Functions Sqrt, LeftLegOf,...
- Variables x, y, a, b,...
- Connectives \neg , \Rightarrow , \land , \lor , \Leftrightarrow
- Equality =
- Quantifiers ∀,∃

Models for FOL: Example



Universal quantification

• ∀<*variables*> <*sentence*>

Everyone at NUS is smart: $\forall x At(x, NUS) \Rightarrow Smart(x)$

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 Roughly speaking, equivalent to the conjunction of instantiations of P

> At(KingJohn,NUS) \Rightarrow Smart(KingJohn) At(Richard,NUS) \Rightarrow Smart(Richard) At(NUS,NUS) \Rightarrow Smart(NUS)

Existential quantification

- ∃<variables> <sentence>
- Someone at NUS is smart:
- $\exists x At(x, NUS) \land Smart(x)$
- Roughly speaking, equivalent to the disjunction of instantiations of P

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At(KingJohn,NUS) ∧ Smart(KingJohn)
∨ At(Richard,NUS) ∧ Smart(Richard)
∨ At(NUS,NUS) ∧ Smart(NUS)
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Properties of quantifiers

- $\forall x \forall y \text{ is the same as } \forall y \forall x$
- $\exists x \exists y \text{ is the same as } \exists y \exists x$

- $\exists x \forall y \text{ is not the same as } \forall y \exists x$
- $\exists x \forall y Loves(x,y)$

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- "There is a person who loves everyone in the world"
- $\forall y \exists x Loves(x,y)$

"Everyone in the world is loved by at least one person"

• Quantifier duality: each can be expressed using the other

Using FOL

The kinship domain:

Brothers are siblings

- $\forall x, y \; Brother(x, y) \Rightarrow Sibling(x, y)$
- One's mother is one's female parent

 $\forall m, c Mother(c) = m \Leftrightarrow (Female(m) \land Parent(m, c))$

One-bit full adder



1. Identify the task

- 2.
- Does the circuit actually add properly? (circuit verification)
- Assemble the relevant knowledge
 3.
 - Composed of wires and gates; Types of gates (AND, OR, XOR, NOT)
 - Irrelevant: size, shape, color, cost of gates

Decide on a vocabulary
 4.

- Encode general knowledge of the domain
 5.
 - $\forall t_1, t_2 \text{ Connected}(t_1, t_2) \Rightarrow \text{Signal}(t_1) = \text{Signal}(t_2)$ $\forall t \text{ Signal}(t) = 1 \lor \text{Signal}(t) = 0$
 - 1 ≠ 0
 - $\forall t_1, t_2 \text{ Connected}(t_1, t_2) \Rightarrow \text{Connected}(t_2, t_1)$
 - $\forall g Type(g) = OR \Rightarrow Signal(Out(1,g)) = 1 \Leftrightarrow \exists n$ Signal(In(n,g)) = 1

Encode general knowledge of the domain
 5.

- → \forall g Type(g) = XOR \Rightarrow Signal(Out(1,g)) = 1 \Leftrightarrow Signal(In(1,g)) ≠ Signal(In(2,g))

 $- \forall g Type(g) = NOT \Rightarrow Signal(Out(1,g)) \neq Signal(In(1,g))$

5. Encode the specific problem instance
 6.

Type $(X_1) = XOR$ Type $(A_1) = AND$ Type $(O_1) = OR$ $Type(X_2) = XOR$ $Type(A_2) = AND$

Connected(Out(1,X₁),In(1,X₂)) Connected(Out(1,X₁),In(2,A₂)) Connected(Out(1,A₂),In(1,O₁)) Connected(Out(1,A₁),In(2,O₁)) Connected(Out(1,X₂),Out(1,C₁)) Connected(Out(1,O₁),Out(2,C₁)) Connected($In(1,C_1),In(1,X_1)$) Connected($In(1,C_1),In(1,A_1)$) Connected($In(2,C_1),In(2,X_1)$) Connected($In(2,C_1),In(2,A_1)$) Connected($In(3,C_1),In(2,X_2)$) Connected($In(3,C_1),In(1,A_2)$)

6. Pose queries to the inference procedure
7.
7. What are the possible sets of values of all the terminals for the adder circuit?

 $\exists i_1, i_2, i_3, o_1, o_2 \text{ Signal}(\text{In}(1, \mathbb{C}_1)) = i_1 \land \text{ Signal}(\text{In}(2, \mathbb{C}_1)) = i_2 \land$ Signal(In(3, \mathbb{C}_1)) = $i_3 \land \text{ Signal}(\text{Out}(1, \mathbb{C}_1)) = o_1 \land$ Signal(Out(2, \mathbb{C}_1)) = o_2

Summary

• First-order logic:

objects and relations are semantic primitives
 syntax: constants, functions, predicates, equality, quantifiers

 Increased expressive power: sufficient to define wumpus world